

# Chapter 2 Categories

## 2.4 representable functor and adjoint functor

- Representable functor
- Adjoint functor

### (I) Representable functor

For any category  $\mathcal{C}$ , there are two functors from  $\mathcal{C}$  to  $\text{Set}$

(1) Covariant  $h_A = \text{Hom}_{\mathcal{C}}(A, \bullet)$

(2) Contravariant  $h^B = \text{Hom}_{\mathcal{C}}(\bullet, B)$

Def. 4.1 Let  $\mathcal{C}$  be a category,  $F: \mathcal{C} \rightarrow \text{Set}$  a covariant functor. If there is  $A \in \text{Ob } \mathcal{C}$  and natural isomorphism

$$\alpha: \text{Hom}_{\mathcal{C}}(A, \bullet) \longrightarrow F$$

we say  $F$  is representable. The pair  $(A, \alpha)$  is called representative of  $F$ .

Similarly, contravariant  $F: \mathcal{C} \rightarrow \text{Set}$  is called representable iff there exists  $B$  such that  $F \cong h^B = \text{Hom}_{\mathcal{C}}(\bullet, B)$ .

Example 4.1 Let  $\mathcal{C}$  be a concrete category, and  $X \in \text{Ob } \mathcal{C}$ , define a set valued functor  $F: \mathcal{C} \rightarrow \text{Set}$  as follows

$$F(A) = \text{Hom}_{\text{Set}}(X, A).$$

$$F(f) = \text{Hom}_{\text{Set}}(X, f) = f_*$$

Suppose  $V$  is a free object over  $X$ , meaning there is  $i: X \rightarrow V$  such that for any  $g: X \rightarrow Y$ , there is a unique  $\tilde{g}$  such that  $\tilde{g} \circ i = g$

$$\begin{array}{ccc} X & \xrightarrow{i} & V \\ & \searrow g & \downarrow \exists! \tilde{g} \\ & & Y \end{array}$$

$F$  is representable with representative  $(V, \alpha)$  where  $\alpha: \text{Hom}(V, \bullet) \rightarrow F$  is defined as

$$\begin{aligned} \alpha_A: \text{Hom}_{\mathcal{C}}(V, A) &\longrightarrow \text{Hom}_{\text{Set}}(X, A) \\ \tilde{g} &\longmapsto \tilde{g} \circ i \end{aligned}$$

Example 4.2. Let  $\text{Mod}_R$  be the category of  $R$  module, and  $A, B \in \text{Mod}_R$ . Define functor  $F: \text{Mod}_R \rightarrow \text{Set}$  as follows:

$$\begin{array}{ccc} C & \longmapsto & \text{Hom}_{\text{bilinear}}(A \times B, C) = F(C) \\ \downarrow g & \rightsquigarrow & F(g) \\ D & & \end{array}$$

$$A \times B \xrightarrow{\text{bilinear}} \begin{array}{c} F(C) \\ \downarrow \\ F(D) \end{array} \xrightarrow{f} C \xrightarrow{g} D$$

Then  $F$  is representable by  $(A \otimes_R B, \alpha)$  with  $\alpha$  defined as

$$\alpha: \text{Hom}_{\text{Mod}_R}(A \otimes_R B, \cdot) \longrightarrow F$$

$$\begin{array}{ccc} \text{Hom}_{\text{Mod}_R}(A \otimes_R B, C) & \xrightarrow{\alpha_C} & F(C) = \text{Hom}_{\text{bilinear}}(A \times B, C) \\ h & \longmapsto & h \circ i \end{array}$$

where  $i: A \times B \longrightarrow A \otimes_R B$  is canonical map.

Lemma (Yoneda lemma). Let  $F: \mathcal{C} \longrightarrow \text{Set}$  be a functor, then the natural transformations from  $F$  to  $h_A = \text{Hom}_{\mathcal{C}}(A, \cdot)$  are in one-to-one correspondence with  $F(A)$ :

$$\text{Nat}(F, h_A) \cong F(A).$$

Moreover, this isomorphism is natural in  $A$  and  $F$  when both sides are regarded as functors from  $\text{Fun}(\mathcal{C}, \text{Set}) \times \mathcal{C}$  to  $\text{Set}$ .

Example.  $\text{Nat}(h_A, h_B) \cong h_A(B) = \text{Hom}_{\mathcal{C}}(A, B)$ .

## (II) Adjoint functor

Def. For two categories  $\mathcal{C}$  and  $\mathcal{D}$ , the product category  $\mathcal{C} \times \mathcal{D}$  is defined as

- $\text{Ob } \mathcal{C} \times \mathcal{D} = \{ (X, Y) \mid X \in \text{Ob } \mathcal{C}, Y \in \text{Ob } \mathcal{D} \}$
- $\text{Hom}_{\mathcal{C} \times \mathcal{D}}((X, Y), (A, B)) = \text{Hom}(X, A) \times \text{Hom}(Y, B)$

$$(X, Y) \xrightarrow{(f, g)} (A, B) \xrightarrow{(k, l)} (C, D)$$

$(k \circ f, l \circ g)$

- We can consider  $F: \mathcal{C} \times \mathcal{D} \rightarrow \text{Set}$ , a typical example is  $\text{Hom}(\cdot, \cdot)$

covariant  
↓  
↑  
contravariant

$$\text{Hom}_{\mathcal{C}}(\cdot, \cdot) : \mathcal{C}^{\text{op}} \times \mathcal{C} \longrightarrow \text{Set}$$

$$(X, Y) \longmapsto \text{Hom}_{\mathcal{C}}(X, Y)$$

$$\begin{array}{ccc} (X, Y) & & \text{Hom}_{\mathcal{C}}(X, Y) & & S \in \text{Hom}_{\mathcal{C}}(X, Y) \\ \downarrow (f, g) & \rightsquigarrow & \downarrow \text{Hom}_{\mathcal{C}}(f, g) & & \downarrow \\ (X', Y') & & \text{Hom}_{\mathcal{C}}(X', Y') & & g \circ S \circ f^{\text{op}} \end{array}$$

$$\begin{array}{ccc} X & \xrightarrow{S} & Y \\ \uparrow f^{\text{op}} & & \downarrow g \\ X' & \longrightarrow & Y' \end{array}$$

- Consider covariant functors  $F: \mathcal{C} \rightarrow \mathcal{D}$ ,  $G: \mathcal{D} \rightarrow \mathcal{C}$ . Then both of

- $\text{Hom}_{\mathcal{D}}(F(\cdot), \cdot)$
- $\text{Hom}_{\mathcal{C}}(\cdot, G(\cdot))$

are functors from  $\mathcal{C}^{\text{op}} \times \mathcal{D}$  to  $\text{Set}$ .

A natural transformation  $\alpha: \text{Hom}_{\mathcal{C}}(\cdot, G(\cdot)) \rightarrow \text{Hom}_{\mathcal{D}}(F(\cdot), \cdot)$  is a set of maps

$$\alpha_{X, Y} : \text{Hom}_{\mathcal{C}}(X, G(Y)) \longrightarrow \text{Hom}_{\mathcal{D}}(F(X), Y)$$

such that for  $f: X \rightarrow X$ ,  $g: Y \rightarrow Y'$ , the following diagram commutes

$$\begin{array}{ccc} \text{Hom}_{\mathcal{C}}(X, G(Y)) & \xrightarrow{\alpha_{X, Y}} & \text{Hom}_{\mathcal{D}}(F(X), Y) \\ \downarrow \text{Hom}_{\mathcal{C}}(f, G(g)) & & \downarrow \text{Hom}_{\mathcal{D}}(F(f), g) \\ \text{Hom}_{\mathcal{C}}(X', G(Y')) & \longrightarrow & \text{Hom}_{\mathcal{D}}(F(X'), Y') \end{array}$$

Def 4.2 Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  and  $G: \mathcal{D} \rightarrow \mathcal{C}$  be two functors. If there exists a natural isomorphism

$$\alpha: \text{Hom}_{\mathcal{C}}(C\bullet, G\bullet) \longrightarrow \text{Hom}_{\mathcal{D}}(F\bullet, \bullet)$$

we say that  $F$  is left adjoint of  $G$  and  $G$  is right adjoint of  $F$ .

Denote  $F \dashv G$ .

Example. In section 7 we prove that

$$\text{Hom}_R(CA \otimes_R B, C) \cong \text{Hom}_R(A, \text{Hom}_R(B, C))$$

This means  $\bullet \otimes_R B \dashv \text{Hom}(B, \bullet)$ .

Prop 4.2 Functor  $G: \mathcal{D} \rightarrow \mathcal{C}$  has left adjoint iff for any  $C \in \text{ob } \mathcal{C}$ , hom-functor  $\text{Hom}(C\bullet, G\bullet)$  is representable.

Proof. " $\Rightarrow$ " Let  $F \dashv G$ , then there is natural isomorphism

$$\alpha_{C, D}: \text{Hom}_{\mathcal{D}}(F(C), D) \xrightarrow{\cong} \text{Hom}_{\mathcal{C}}(C, G(D))$$

Fix  $C$ , we see  $\text{Hom}_{\mathcal{D}}(F(C), \bullet) \cong \text{Hom}_{\mathcal{C}}(C, G(\bullet))$ , meaning  $\text{Hom}(C, G(\bullet))$  is representable.

" $\Leftarrow$ " Suppose  $(A\bullet, \alpha)$  be representative of  $\text{Hom}_{\mathcal{C}}(C, G(\bullet))$

Define  $F(C) = A\bullet$ , we can check

$$\text{Hom}(F(\bullet), \bullet) \xrightarrow{\cong} \text{Hom}(\bullet, G(\bullet)).$$